TRUE
$$\triangleq \lambda x \cdot (\lambda y \cdot x)$$

FALSE $\triangleq \lambda x \cdot \lambda y \cdot y$
 $IF \triangleq \lambda b \cdot \lambda t \cdot \lambda f \cdot b$
 $b t f$

$$(\lambda \times \times \times) (\lambda \times \times \times)$$

$$e ::= x | e_1 e_2 | \lambda x : \tau . e |$$

$$n | true | false$$

$$\tau ::= bool | int | \gamma_1 \to \gamma_2$$

$$bool + int$$

$$(int \rightarrow bool) \rightarrow int$$

$$\begin{cases} y \cdot ball \\ + 4 + \lambda &: int \\ + \lambda \times (int, x) \quad int \\ + \lambda \times (int, x) \quad H : int \\ f(\lambda, \cdot, int, x) \quad H : int \\ f(\lambda, \cdot, int, x) \quad H : int \\ f(\lambda, \cdot, int, x) \quad H : int \\ \hline f(\lambda, \cdot, int, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\ \hline f(\lambda, \cdot, x) \quad H : int \\$$

 $F(\lambda \times 100) - \times +38) H : int$

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WELL-TYPE PRESERVE CANNOT
CONDANCESS
IF E IS MELL-TYPED AND E * 5
AND E' IS IRREDUCIBLE, THEN E'
AND E' IS IRREDUCIBLE, THEN E'
A VANE.

$$A = (1+1) \rightarrow 1+2 \rightarrow 6$$

 $(A_X: int. x) + H$
 $H + (1+1) \rightarrow 1+2 \rightarrow 6$
 $(A_X: int. x) + H$
 $H + (1+1) \rightarrow 1+2 \rightarrow 6$
 $(A_X: int. x) + H$
 $H + (1+1) \rightarrow 1+2 \rightarrow 6$
 $(X: int. x)$
 $M = (A_X, X \times) (A_X, X \times)$
 $\omega = (A_X, X \times) (A_X, X \times)$
 $\omega = (A_X, X \times) (A_X, X \times)$
 $\omega = (A_X, X + (A_X, X \times)) (A_X, X \times)$
 $\omega = (A_X, X + (A_X, X \times)) (A_X, X \times)$
 $\chi = \chi + \chi'$
 $\chi = \chi + \chi'$

$$\frac{\Gamma + e_{i}\gamma_{i}\Gamma + e_{2}\gamma_{2}}{\Gamma + (e_{i}, e_{2})\gamma_{i}\gamma_{i}\gamma_{2}} \qquad \frac{\Gamma + e_{i}\gamma_{i}\gamma_{2}}{\Gamma + \# Ie^{i}\gamma_{i}\gamma_{1}}$$

$$\frac{\left[\frac{1}{1 + e} \cdot \frac{1}{2} \times \frac{1}{2} \right]}{\left[\frac{1}{1 + 2} \cdot \frac{1}{2} \cdot \frac{1}{2} \right]}$$

$$e ::= \cdots \left[\frac{1}{1 + 2} \cdot \frac{1}{2} \right]$$

$$watch e with e_1 e_2$$

$$V ::= \cdots \left[\frac{1}{2} \cdot \frac{1}{1 + 2} \right]$$

$$\frac{1}{1 + 2}$$

$$V ::= \cdots \left[\frac{1}{2} \cdot \frac{1}{1 + 2} \right]$$

$$int$$

$$int$$

$$int$$

$$int$$

$$int + bool$$

$$(int int + bool true) : int + bool$$

$$(int int + bool true) : int + bool$$

$$(Ax : int. 0)$$

$$int$$

$$\frac{\Gamma + e \cdot \tau_{1}}{\Gamma + i \alpha L_{\alpha_{1} + \tau_{2}} e \cdot \tau_{1} + \tau_{2}} \qquad \frac{\Gamma + e \cdot \tau_{2}}{\Gamma + i \alpha L_{\alpha_{1} + \tau_{2}} e \cdot \tau_{1} + \tau_{2}} \qquad \frac{\Gamma + e \cdot \tau_{1} + \tau_{2}}{\Gamma + e_{1} \cdot \tau_{1} + \tau_{3}} \qquad \frac{\Gamma + e_{2} \cdot \tau_{2} + \tau_{3}}{\Gamma + i \alpha t c h e with e \cdot le_{2} \cdot \tau_{3}}$$

$$\lambda x : int \cdot x$$

 $\lambda x : bool \cdot x$
 $\lambda x : (int \to bool) \cdot x$

 $\left(\bigwedge \alpha \cdot \chi \times \alpha \cdot \chi \right) (int \rightarrow bool)$ $\rightarrow \chi \times \frac{(int)}{bool} \times$ $\bigwedge \alpha \cdot \bigwedge \beta \cdot \chi \times \beta \cdot \beta \cdot (\chi, \chi)$